## Computer Algebra with D-Finite Functions

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June 20, 2024<br>Radboud University Nijmegen

ÖAW RICAM

## Special Functions

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- arise in mathematical analysis and in real-world phenomena
- are solutions to certain differential equations
- cannot be expressed in terms of the usual elementary functions $(\sqrt{ }, \exp , \log , \sin , \cos , \ldots)$


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## D-finite Functions

Definition: A function $f(x)$ is called D-finite if it satisfies a linear ordinary differential equation with polynomial coefficients:

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p_{r}(x) f^{(r)}(x)+\cdots+p_{1}(x) f^{\prime}(x)+p_{0}(x) f(x)=0
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$p_{0}, \ldots, p_{r} \in \mathbb{K}[x]$ (not all zero).

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- $a_{n}$ is P-recursive if and only if $f(x)=\sum_{n=0}^{\infty} a_{n} x^{n}$ is D-finite.


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- Equivalently, such functions/sequences are called holonomic.


## Differential Equations and Recurrences

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Many special functions can be characterized as solutions to systems of linear differential equations and recurrences, and in fact are D-finite (holonomic).

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\left(x^{2}-1\right) P_{n}^{\prime \prime}(x)+2 x P_{n}^{\prime}(x)-n(n+1) P_{n}(x)=0 .
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Consider the set $\left\{P_{n}^{(i)}(x): i \geqslant 0\right\}$.

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\begin{aligned}
& P_{n}^{(4)}(x)= \\
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$\longrightarrow P_{n}(x)$ is $\mathbf{D}$-finite w.r.t. $x$.
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& P_{n}^{(4)}(x)= \\
& -\frac{8 x\left(n^{2} x^{2}-n^{2}+n x^{2}-n+3 x^{2}+3\right)}{\left(x^{2}-1\right)^{3}} P_{n}^{\prime}(x) \\
& +\frac{n(n+1)\left(n^{2} x^{2}-n^{2}+n x^{2}-n+18 x^{2}+6\right)}{\left(x^{2}-1\right)^{3}} P_{n}(x)
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Example: The Legendre polynomials can be defined recursively:

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\begin{aligned}
P_{0}(x) & =1, \quad P_{1}(x)=x \\
n P_{n}(x) & =(2 n-1) x P_{n-1}(x)-(n-1) P_{n-2}(x)
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$\longrightarrow P_{n}(x)$ is D-finite w.r.t. $n$ and $x$ (of rank 2).
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- $\int f_{n}(x) \mathrm{d} x$ and $\sum_{n} f_{n}(x)$
- $\frac{\mathrm{d}}{\mathrm{d} x} f_{n}(x)$
- $f_{a n+b}(x)$, where $a, b \in \mathbb{Z}$
- $f_{n}(h(x))$, where $h(x)$ is an algebraic function


## Many Functions are D-finite

ArcCsc, KelvinBei, HypergeometricPFQ, ExpIntegralE, ArcTanh, HankelH2, AngerJ, JacobiP, ChebyshevT, AiryBi, AiryAi, Sinc, Multinomial, CatalanNumber, QBinomial, CosIntegral, ArcSech, SphericalHankelH2, HermiteH, ExplntegralEi, Beta, AiryBiPrime, SphericalBesselJ, Binomial, ParabolicCylinderD, Erfc, EllipticK, Fibonacci, QFactorial, Cos, Hypergeometric2F1, Erf, KelvinKer, HypergeometricPFQRegularized, Log, Factorial, BesselY, Cosh, CoshIntegral, ArcTan, ArcCoth, LegendreP, LaguerreL, EllipticE, SinhIntegral, Sinh, BetaRegularized, SphericalHankelH1, ArcSin, EllipticThetaPrime, Root, LucasL, AppellF1, FresneIC, LegendreQ, ChebyshevU, GammaRegularized, Erfi, HarmonicNumber, Bessell, KelvinKei, ArithmeticGeometricMean, Exp, ArcCot, EllipticTheta, Hypergeometric0F1, EllipticPi, GegenbauerC, ArcCos, WeberE, FresneIS, EllipticF, ArcCosh, Subfactorial, QPochhammer, Gamma, StruveH, WhittakerM, ArcCsch, Hypergeometric1F1, SinIntegral, BesselJ, StruveL, ArcSec, Factorial2, KelvinBer, BesselK, ArcSinh, HankelH1, Sqrt, PolyGamma, HypergeometricU, AiryAiPrime, Sin,

## Quiz: Who is D-Finite?

$$
\operatorname{erf}(\sqrt{x+1})^{2}+\exp (\sqrt{x+1})^{2}
$$

$$
\left((\sinh (x))^{2}+(\sin (x))^{-2}\right) \cdot\left((\cosh (x))^{2}+(\cos (x))^{-2}\right)
$$

$$
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## Quiz: Who is D-Finite?

$$
\operatorname{erf}(\sqrt{x+1})^{2}+\exp (\sqrt{x+1})^{2}
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## The Symbolic Computation Viewpoint

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The holonomic systems approach (Zeilberger 1990) is a versatile toolbox for solving many different kinds of mathematical problems:

- calculate integrals and summation formulas
- prove special function identities
- computations in $q$-calculus (e.g., quantum knot invariants)
- fast numerical evaluation of mathematical functions
- number theory (e.g., irrationality proofs)
- evaluate symbolic determinants (e.g., in combinatorics)


## Application

## Finite Elements


(joint work with Joachim Schöberl and Peter Paule)

## Problem Setting

Simulate the propagation of electromagnetic waves according to

$$
\begin{equation*}
\frac{\mathrm{d} H}{\mathrm{~d} t}=\operatorname{curl} E, \quad \frac{\mathrm{~d} E}{\mathrm{~d} t}=-\operatorname{curl} H \tag{Maxwell}
\end{equation*}
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where $H$ and $E$ are the magnetic and the electric field respectively.

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Define basis functions (2D case):

$$
\varphi_{i, j}(x, y):=(1-x)^{i} P_{j}^{(2 i+1,0)}(2 x-1) P_{i}\left(\frac{2 y}{1-x}-1\right)
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using Legendre and Jacobi polynomials.

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using Legendre and Jacobi polynomials.

Problem: Represent the partial derivatives of $\varphi_{i, j}(x, y)$ in the basis (i.e., as linear combinations of shifts of the $\varphi_{i, j}(x, y)$ itself).

## Solution

Ansatz: One needs a relation of the form

$$
\sum_{(k, l) \in A} a_{k, l}(i, j) \frac{\mathrm{d}}{\mathrm{~d} x} \varphi_{i+k, j+l}(x, y)=\sum_{(m, n) \in B} b_{m, n}(i, j) \varphi_{i+m, j+n}(x, y)
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that is free of $x$ and $y$ (and similarly for $\frac{\mathrm{d}}{\mathrm{d} y}$ ).

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that is free of $x$ and $y$ (and similarly for $\frac{\mathrm{d}}{\mathrm{d} y}$ ).
Result: Computer algebra methods (D-finite closure properties, Gröbner bases), deliver the relation

$$
\begin{aligned}
& (2 i+j+3)(2 i+2 j+7) \frac{\mathrm{d}}{\mathrm{~d} x} \varphi_{i, j+1}(x, y)+ \\
& 2(2 i+1)(i+j+3) \frac{\mathrm{d}}{\mathrm{~d} x} \varphi_{i, j+2}(x, y)- \\
& (j+3)(2 i+2 j+5) \frac{\mathrm{d}}{\mathrm{~d} x} \varphi_{i, j+3}(x, y)+ \\
& (j+1)(2 i+2 j+7) \frac{\mathrm{d}}{\mathrm{~d} x} \varphi_{i+1, j}(x, y)- \\
& 2(2 i+3)(i+j+3) \frac{\mathrm{d}}{\mathrm{~d} x} \varphi_{i+1, j+1}(x, y)- \\
& (2 i+j+5)(2 i+2 j+5) \frac{\mathrm{d}}{\mathrm{~d} x} \varphi_{i+1, j+2}(x, y)+ \\
& 2(i+j+3)(2 i+2 j+5)(2 i+2 j+7) \varphi_{i, j+2}(x, y)+ \\
& 2(i+j+3)(2 i+2 j+5)(2 i+2 j+7) \varphi_{i+1, j+1}(x, y)=0 .
\end{aligned}
$$

## Creative Telescoping



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Creative telescoping is a method

- to deal with parametrized symbolic sums and integrals


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\sum_{k=1}^{\infty} \frac{1}{k(k+n)}
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Example:

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\sum_{k=1}^{\infty} \frac{1}{k^{2}} & =\frac{\pi^{2}}{6} \quad \text { Bad: no parameter! } \\
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\end{aligned}
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Example:

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\begin{gathered}
\sum_{k=1}^{\infty} \frac{1}{k^{2}}=\frac{\pi^{2}}{6} \quad \text { Bad: no parameter! } \\
\underbrace{\sum_{k=1}^{\infty} \frac{1}{k(k+n)}}_{=: f_{n}} \rightsquigarrow(n+2)^{2} f_{n+2}=(n+1)(2 n+3) f_{n+1}-n(n+1) f_{n}
\end{gathered}
$$

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Method for doing integrals and sums (aka Feynman's differentiating under the integral sign)
Consider the following summation problem: $F(n):=\sum_{k=a}^{b} f(n, k)$

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c_{r}(n) f(n+r, k)+\cdots+c_{0}(n) f(n, k)=g(n, k+1)-g(n, k)
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Summing from $a$ to $b$ yields a recurrence for $F(n)$ :

$$
c_{r}(n) F(n+r)+\cdots+c_{0}(n) F(n)=g(n, b+1)-g(n, a)
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## Creative Telescoping

Method for doing integrals and sums (aka Feynman's differentiating under the integral sign)
Consider the following integration problem: $F(x):=\int_{a}^{b} f(x, y) \mathrm{d} y$
Telescoping: write $f(x, y)=\frac{\mathrm{d}}{\mathrm{d} y} g(x, y)$.
Then $F(n)=\int_{a}^{b}\left(\frac{\mathrm{~d}}{\mathrm{~d} y} g(x, y)\right) \mathrm{d} y \quad=g(x, b)-g(x, a)$.
Creative Telescoping: write

$$
c_{r}(x) \frac{\mathrm{d}^{r}}{\mathrm{~d} x^{r}} f(x, y)+\cdots+c_{0}(x) f(x, y)=\frac{\mathrm{d}}{\mathrm{~d} y} g(x, y)
$$

Integrating from $a$ to $b$ yields a differential equation for $F(x)$ :

$$
c_{r}(x) \frac{\mathrm{d}^{r}}{\mathrm{~d} x^{r}} F(x)+\cdots+c_{0}(x) F(x)=g(x, b)-g(x, a)
$$

## Application

## Special Function Identities



## Table of Integrals by Gradshteyn and Ryzhik

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$\qquad$


I. S. GRADSHTEYN

I. M. RYZHIK
8
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INTEGRALS, SERIES, AND PRODUCTS
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## Table of Integrals by Gradshteyn and Ryzhik

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## Table of Integrals by Gradshteyn and Ryzhik

1. 

$$
\begin{aligned}
& \text { 1. } \begin{array}{l}
\int_{0}^{1}(1-x)^{\mu-1} x^{\nu-1} C_{2 n}^{\lambda}\left(\gamma x^{1 / 2}\right) d x=(-1)^{n} \frac{\Gamma(\lambda+n) \Gamma(\mu) \Gamma(\nu)}{n!\Gamma(\lambda) \Gamma(\mu+\nu)}{ }_{3} F_{2}\left(-n, n+\lambda, \nu ; \frac{1}{2}, \mu+\nu ; \gamma^{2}\right) \\
{[\operatorname{Re} \mu>0, \quad \operatorname{Re} \nu>0] \quad \text { ET II 191(41)a }} \\
2 . \quad \int_{0}^{1}(1-x)^{\mu-1} x^{\nu-1} C_{2 n+1}^{\lambda}\left(\gamma x^{1 / 2}\right) d x=\frac{(-1)^{n} 2 \gamma \Gamma(\mu) \Gamma(\lambda+n+1) \Gamma\left(\nu+\frac{1}{2}\right)}{n!\Gamma(\lambda) \Gamma\left(\mu+\nu+\frac{1}{2}\right)} \\
\times{ }_{3} F_{2}\left(-n, n+\lambda+1, \nu+\frac{1}{2} ; \frac{3}{2}, \mu+\nu+\frac{1}{2} ; \gamma^{2}\right) \\
{\left[\operatorname{Re} \mu>0, \quad \operatorname{Re} \nu>-\frac{1}{2}\right] \quad \text { ET II 191(42) }}
\end{array}
\end{aligned}
$$

### 7.32 Combinations of Gegenbauer polynomials $C_{n}^{\nu}(x)$ and elementary functions

 7.321$$
\begin{array}{r}
\int_{-1}^{1}\left(1-x^{2}\right)^{\nu-\frac{1}{2}} e^{i a x} C_{n}^{\nu}(x) d x=\frac{\pi 2^{1-\nu} i^{n} \Gamma(2 \nu+n)}{n!\Gamma(\nu)} a^{-\nu} J_{\nu+n}(a) \\
{\left[\operatorname{Re} \nu>-\frac{1}{2}\right]}
\end{array}
$$

ET II 281(7), MO 99a
7.322

$$
\int_{0}^{2 a}[x(2 a-x)]^{\nu-\frac{1}{2}} C_{n}^{\nu}\left(\frac{x}{a}-1\right) e^{-b x} d x=(-1)^{n} \frac{\pi \Gamma(2 \nu+n)}{n!\Gamma(\nu)}\left(\frac{a}{2 b}\right)^{\nu} e^{-a b} I_{\nu+n}(a b)
$$

$$
\left[\operatorname{Re} \nu>-\frac{1}{2}\right]
$$

ET I 171(9)
7.323
1.
$\int_{0}^{\pi} C_{n}^{\nu}(\cos \varphi)(\sin \varphi)^{2 \nu} d \varphi=0$

$$
[n=1,2,3, \ldots]
$$

## Table of Integrals by Gradshteyn and Ryzhik

$$
\int_{-1}^{1}\left(1-x^{2}\right)^{\nu-\frac{1}{2}} e^{i a x} C_{n}^{\nu}(x) d x=\frac{\pi 2^{1-\nu} i^{n} \Gamma(2 \nu+n)}{n!\Gamma(\nu)} a^{-\nu} J_{\nu+n}(a)
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Gegenbauer
polynomials $C_{n}^{(\alpha)}(x)$

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Gegenbauer polynomials $C_{n}^{(\alpha)}(x)$

Gamma
function $\Gamma(x)$


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- A large portion of such identities can be proven via the holonomic systems approach.


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Gamma
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- A large portion of such identities can be proven via the holonomic systems approach.
- Algorithms are implemented in the HolonomicFunctions package.


## The HolonomicFunctions Package

Example: Holonomic system, satisfied by both sides of the identity:

$$
\begin{aligned}
& i a(n+2 \nu) f_{n}^{\prime}(a)+a(n+1) f_{n+1}(a)-i n(n+2 \nu) f_{n}(a)=0 \\
& a(n+1)(n+2) f_{n+2}(a)-2 i(n+1)(n+\nu+1)(n+2 \nu+1) f_{n+1}(a) \\
& \quad-a(n+2 \nu)(n+2 \nu+1) f_{n}(a)=0
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\end{aligned}
$$

$\operatorname{In}[42]:=$ Annihilator [Pi * $2^{\wedge}(1-v) * I^{\wedge} n * \operatorname{Gamma}[2 v+n] / n!/ G a m m a[v] * a^{\wedge}(-v)$ * BesselJ[v + n, a], \{Der[a], S[n]\}] // Factor

Out[42]=

$$
\begin{aligned}
& \left\{\text { ii } a(n+2 v) D_{a}+a(1+n) S_{n}-i \operatorname{n}(\mathrm{n}+2 v),\right. \\
& \left.\mathrm{a}(1+\mathrm{n})(2+\mathrm{n}) \mathrm{S}_{\mathrm{n}}^{2}-2 \text { ii }(1+\mathrm{n})(1+\mathrm{n}+v)(1+\mathrm{n}+2 v) \mathrm{S}_{\mathrm{n}}-\mathrm{a}(\mathrm{n}+2 v)(1+\mathrm{n}+2 v)\right\}
\end{aligned}
$$

$\operatorname{In}[43]:=$ CreativeTelescoping $\left[\left(1-x^{\wedge} 2\right)^{\wedge}(v-1 / 2) * \operatorname{Exp}[I * a * x] * \operatorname{GegenbauerC}[n, v, x]\right.$, $\operatorname{Der}[\mathrm{x}],\{\operatorname{Der}[\mathrm{a}], \mathrm{S}[\mathrm{n}]\}] / /$ Factor
Out[43]=

$$
\begin{aligned}
\{ & \left\{a(n+2 v) D_{a}-i \operatorname{a}(1+n) S_{n}-n(n+2 v),\right. \\
& \left.a(1+n)(2+n) S_{n}^{2}-2 i(1+n)(1+n+v)(1+n+2 v) S_{n}-a(n+2 v)(1+n+2 v)\right\} \\
& \left.\left\{(1+n) S_{n}-x(n+2 v), 2 i(1+n) x(1+n+v) S_{n}-2 i(1+n+v)(n+2 v)\right\}\right\}
\end{aligned}
$$

## Special Function Identities

$$
\begin{equation*}
\sum_{k=0}^{n}\binom{n}{k}^{2}\binom{k+n}{k}^{2}=\sum_{k=0}^{n}\binom{n}{k}\binom{k+n}{k} \sum_{j=0}^{k}\binom{k}{j}^{3} \tag{1}
\end{equation*}
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& \int_{0}^{\infty} \frac{1}{\left(x^{4}+2 a x^{2}+1\right)^{m+1}} \mathrm{~d} x=\frac{\pi P_{m}^{\left(m+\frac{1}{2},-m-\frac{1}{2}\right)}(a)}{2^{m+\frac{3}{2}}(a+1)^{m+\frac{1}{2}}} \tag{2}
\end{align*}
$$

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\int_{0}^{\infty} \frac{1}{\left(x^{4}+2 a x^{2}+1\right)^{m+1}} \mathrm{~d} x=\frac{\pi P_{m}^{\left(m+\frac{1}{2},-m-\frac{1}{2}\right)}(a)}{2^{m+\frac{3}{2}}(a+1)^{m+\frac{1}{2}}}  \tag{2}\\
e^{-x} x^{a / 2} n!L_{n}^{a}(x)=\int_{0}^{\infty} e^{-t} t^{\frac{a}{2}+n} J_{a}(2 \sqrt{t x}) \mathrm{d} t \tag{3}
\end{gather*}
$$

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\begin{align*}
& \sum_{k=0}^{n}\binom{n}{k}^{2}\binom{k+n}{k}^{2}=\sum_{k=0}^{n}\binom{n}{k}\binom{k+n}{k} \sum_{j=0}^{k}\binom{k}{j}^{3}  \tag{1}\\
& \int_{0}^{\infty} \frac{1}{\left(x^{4}+2 a x^{2}+1\right)^{m+1}} \mathrm{~d} x=\frac{\pi P_{m}^{\left(m+\frac{1}{2},-m-\frac{1}{2}\right)}(a)}{2^{m+\frac{3}{2}}(a+1)^{m+\frac{1}{2}}}  \tag{2}\\
& e^{-x} x^{a / 2} n!L_{n}^{a}(x)=\int_{0}^{\infty} e^{-t} t^{\frac{a}{2}+n} J_{a}(2 \sqrt{t x}) \mathrm{d} t  \tag{3}\\
& \int_{-\infty}^{\infty} \sum_{m=0}^{\infty} \sum_{n=0}^{\infty} \frac{H_{m}(x) H_{n}(x) r^{m} s^{n} e^{-x^{2}}}{m!n!} \mathrm{d} x=\sqrt{\pi} e^{2 r s} \tag{4}
\end{align*}
$$

## Special Function Identities

$$
\begin{gather*}
\sum_{k=0}^{n}\binom{n}{k}^{2}\binom{k+n}{k}^{2}=\sum_{k=0}^{n}\binom{n}{k}\binom{k+n}{k} \sum_{j=0}^{k}\binom{k}{j}^{3}  \tag{1}\\
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\int_{-\infty}^{\infty} \sum_{m=0}^{\infty} \sum_{n=0}^{\infty} \frac{H_{m}(x) H_{n}(x) r^{m} s^{n} e^{-x^{2}}}{m!n!} \mathrm{d} x=\sqrt{\pi} e^{2 r s}  \tag{4}\\
\int_{-1}^{1}\left(1-x^{2}\right)^{\nu-\frac{1}{2}} e^{i a x} C_{n}^{(\nu)}(x) \mathrm{d} x=\frac{\pi i^{n} \Gamma(n+2 \nu) J_{n+\nu}(a)}{2^{\nu-1} a^{\nu} n!\Gamma(\nu)} \tag{5}
\end{gather*}
$$

## Application

## Determinants Count!

(joint work with Hao Du,
Christian Krattenthaler, Michael Schlosser, Aek Thanatipanonda, Elaine Wong)


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## Symbolic Determinants via Holonomic Ansatz

$$
\operatorname{det}_{1 \leqslant i, j \leqslant n} \frac{1}{i+j-1}=\frac{1}{(2 n-1)!} \prod_{k=1}^{n-1} \frac{(k!)^{2}}{(k+1)_{n-1}}
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\operatorname{det}_{0 \leqslant i, j \leqslant n-1}\binom{2 i+2 a}{j+b} & =2^{n(n-1) / 2} \prod_{k=0}^{n-1} \frac{(2 k+2 a)!k!}{(k+b)!(2 k+2 a-b)!}
\end{aligned}
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\operatorname{det}_{0 \leqslant i, j \leqslant n-1} \sum_{k}\binom{i}{k}\binom{j}{k} 2^{k} & =2^{n(n-1) / 2}
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\operatorname{det}_{0 \leqslant i, j \leqslant n-1} \sum_{k}\binom{i}{k}\binom{j}{k} 2^{k}=2^{n(n-1) / 2} \\
1 \leqslant i, j \leqslant 2 m+1 \\
=\frac{(-1)^{m-r+1}(\mu+3)(m+r+1)_{m-r}}{\operatorname{det}^{m}} \cdot \prod_{i=1}^{2 m} \frac{(\mu+i+3)_{2 r}}{(i)_{2 r}} \\
\\
\times \prod_{i=1}^{m-2 r+1} \frac{\left(\frac{\mu}{2}+r+\frac{3}{2}\right)_{m-r+1}}{} \frac{(\mu+2 i+6 r+3)_{i}^{2}\left(\frac{\mu}{2}+2 i+3 r+2\right)_{i-1}^{2}}{(i)_{i}^{2}\left(\frac{\mu}{2}+i+3 r+2\right)_{i-1}^{2}}
\end{gathered}
$$

The Holonomic Ansatz

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## The Holonomic Ansatz

Problem: Prove a determinantal identity of the form $\operatorname{det}_{1 \leqslant i, j \leqslant n}\left(a_{i, j}\right)=b_{n}$


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Problem: Prove a determinantal identity of the form $\operatorname{det}_{1 \leqslant i, j \leqslant n}\left(a_{i, j}\right)=b_{n}$, where

- $a_{i, j}$ is a holonomic sequence
- that does not depend on $n$



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Problem: Prove a determinantal identity of the form $\operatorname{det}_{1 \leqslant i, j \leqslant n}\left(a_{i, j}\right)=b_{n}$, where

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- that does not depend on $n$, and
- $b_{n}$ is a closed form $\left(b_{n} \neq 0\right.$ for all $\left.n\right)$.



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$$
\mathcal{A}_{n}=\left(\begin{array}{c:c} 
& \\
\mathcal{A}_{n-1} & \\
& \\
\hdashline a_{n, 1} & \cdots
\end{array} a_{n, n-1}: a_{n, n}\right)
$$



Laplace expansion:

$$
\operatorname{det}\left(\mathcal{A}_{n}\right)=a_{n, 1} \operatorname{Cof}_{n, 1}+\cdots+a_{n, n-1} \operatorname{Cof}_{n, n-1}+a_{n, n} \operatorname{det}\left(\mathcal{A}_{n-1}\right)
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$\sum(n)^{2}\binom{3 m k}{2 n}=\binom{n n}{n}^{2}$ WHO YOU GONNA CALL?

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\frac{\operatorname{det}\left(\mathcal{A}_{n}\right)}{\operatorname{det}\left(\mathcal{A}_{n-1}\right)}=a_{n, 1} \frac{\operatorname{Cof}_{n, 1}}{\operatorname{det}\left(\mathcal{A}_{n-1}\right)}+\cdots+a_{n, n-1} \frac{\operatorname{Cof}_{n, n-1}}{\operatorname{det}\left(\mathcal{A}_{n-1}\right)}+a_{n, n}
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Laplace expansion:

$$
0=\sum_{j=1}^{n} a_{i, j} c_{n, j} \quad(1 \leqslant i<n), \quad c_{n, n}=1
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## Recipe

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c_{n, n} & =1 & & (1 \leqslant n)  \tag{1}\\
\sum_{j=1}^{n} a_{i, j} c_{n, j} & =0 & & (1 \leqslant i<n)  \tag{2}\\
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Conjecture (Di Francesco's determinant for 20V configurations):

$$
\operatorname{det}_{0 \leqslant i, j<n}\left(2^{i}\binom{i+2 j+1}{2 j+1}-\binom{i-1}{2 j+1}\right)=2 \prod_{i=1}^{n} \frac{2^{i-1}(4 i-2)!}{(n+2 i-1)!}
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Theorem (Di Francesco's determinant for 20V configurations):

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$$

## Inverse Inequalities <br> (join work with M. Neumüller and S. Radu)

We consider inequalities of the form

$$
\left\|v_{n}\right\|_{X(\Omega)} \leqslant c(h, n)\left\|v_{n}\right\|_{Y(\Omega)} \quad \text { for all } v_{n} \in V_{n}
$$

- $\Omega \subset \mathbb{R}^{d}, d \in \mathbb{N}$
- $V$ : some infinite-dimensional space of functions defined on $\Omega$
- $\|\cdot\|_{X(\Omega)},\|\cdot\| \|_{Y(\Omega)}$ : norms that are used in numerical methods
- $\left(V_{n}\right)_{n \in \mathbb{N}}$ : finite-dimensional approximation of $V$
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Transform the problem to a reference element $\hat{\Omega}$ :

$$
\hat{c}(n)=\sup _{v_{n} \in \hat{V}_{n}} \frac{\left\|v_{n}\right\|_{X(\hat{\Omega})}}{\left\|v_{n}\right\|_{Y(\hat{\Omega})}}=\sqrt{\sup _{v_{n} \in \hat{V}_{n}} \frac{\left(v_{n}, v_{n}\right)_{X(\hat{\Omega})}}{\left(v_{n}, v_{n}\right)_{Y(\hat{\Omega})}}}
$$

## Inverse Inequalities

Here we consider the reference domain $\hat{\Omega}=(-1,1)^{2}$ with

$$
\begin{aligned}
& (u, v)_{X(\hat{\Omega})}=\int_{\hat{\Omega}} \partial_{x} u(x, y) \partial_{x} v(x, y) \mathrm{d} x \mathrm{~d} y \\
& (u, v)_{Y(\hat{\Omega})}=\int_{\hat{\Omega}} u(x, y) v(x, y) \mathrm{d} x \mathrm{~d} y
\end{aligned}
$$

for $u, v \in \hat{V}_{n}$, where $\hat{V}_{n}$ is the space of polynomials of degree less than $n$, i.e.

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\hat{V}_{n}=\left\{x^{i} y^{j}: 0 \leqslant i, j<n\right\} .
$$

The desired "constant" $\hat{c}(n)$ can be found as the largest $\lambda_{n}$ solving the generalized eigenvalue problem

$$
B_{n} \vec{x}_{n}=\lambda_{n} A_{n} \vec{x}_{n},
$$

where $A_{n}$ and $B_{n}$ are certain $n \times n$ matrices.

## The Matrices

$$
a_{i, j}:=\frac{1-(-1)^{i+j-1}}{i+j-1}, \quad b_{i, j}:=(i-1)(j-1) \frac{1-(-1)^{i+j-3}}{i+j-3}
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\left|B_{6}-\lambda A_{6}\right|=\left|\begin{array}{cccccc}
-2 \lambda & 0 & -\frac{2}{3} \lambda & 0 & -\frac{2}{5} \lambda & 0 \\
0 & 2-\frac{2}{3} \lambda & 0 & 2-\frac{2}{5} \lambda & 0 & 2-\frac{2}{7} \lambda \\
-\frac{2}{3} \lambda & 0 & \frac{8}{3}-\frac{2}{5} \lambda & 0 & \frac{16}{5}-\frac{2}{7} \lambda & 0 \\
0 & 2-\frac{2}{5} \lambda & 0 & \frac{18}{5}-\frac{2}{7} \lambda & 0 & \frac{30}{7}-\frac{2}{9} \lambda \\
-\frac{2}{5} \lambda & 0 & \frac{16}{5}-\frac{2}{7} \lambda & 0 & \frac{32}{7}-\frac{2}{9} \lambda & 0 \\
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\end{array}\right|
\end{gathered}
$$

Hence we get: $\quad \operatorname{det}\left(B_{n}-\lambda A_{n}\right)=2^{n} \operatorname{det}\left(A_{\lceil n / 2\rceil}^{(1)}\right) \cdot \operatorname{det}\left(A_{\lfloor n / 2\rfloor}^{(0)}\right)$.

## The Determinant

By a variation of the holonomic ansatz we prove:
Theorem.

$$
\begin{aligned}
& \operatorname{det} A_{n}^{(0)}=\underbrace{\left(-\frac{1}{2}\right)^{n} \prod_{i=1}^{n} \frac{((i-1)!)^{2}}{\left(i+\frac{1}{2}\right)_{n}}}_{\text {"hyperholonomic" part }} \underbrace{\sum_{j=0}^{n}(-4)^{j-n} \frac{(2 n-2 j+1)_{2 n}}{(2 j)!} \lambda^{j}}_{\text {holonomic part }}, \\
& \operatorname{det} A_{n}^{(1)}=\underbrace{\left(-\frac{1}{2}\right)^{n} \prod_{i=1}^{n} \frac{((i-1)!)^{2}}{\left(i-1+\frac{1}{2}\right)_{n}}}_{\text {"hyperholonomic" part }} \underbrace{\sum_{j=0}^{n-1} \frac{(2 n-2 j-1)_{2 n-1}}{(-4)^{n-j-1}(2 j+1)!} \lambda^{j}}_{\text {holonomic part }} .
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\end{aligned}
$$

We use this explicit evaluation to estimate the largest eigenvalue.

## Final Result

Theorem: For all $n \in \mathbb{N}$ we have the estimate $b_{1}(n)<\lambda_{n}<b_{2}(n)$

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& b_{1}(n):=\frac{m_{1}(n)}{2}\left(1+\sqrt{1-\frac{2}{3} \frac{(n-2)(n-3)(n+3)(n+4)}{n(n-1)(n+1)(n+2)}}\right) \\
& b_{2}(n):=m_{1}(n)\left(\frac{1}{3}+\left(r_{1}(n)+\sqrt{r_{2}(n)}\right)^{1 / 3}+\left(r_{1}(n)-\sqrt{r_{2}(n)}\right)^{1 / 3}\right)
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\begin{aligned}
& b_{1}(n):=\frac{m_{1}(n)}{2}\left(1+\sqrt{1-\frac{2}{3} \frac{(n-2)(n-3)(n+3)(n+4)}{n(n-1)(n+1)(n+2)}}\right) \\
& b_{2}(n):=m_{1}(n)\left(\frac{1}{3}+\left(r_{1}(n)+\sqrt{r_{2}(n)}\right)^{1 / 3}+\left(r_{1}(n)-\sqrt{r_{2}(n)}\right)^{1 / 3}\right)
\end{aligned}
$$

where $m_{1}, r_{1}$, and $r_{2}$ are given by

$$
\begin{aligned}
m_{1}(n) & :=\frac{n(n-1)(n+1)(n+2)}{8} \\
r_{1}(n) & :=\frac{2\left(n^{8}+4 n^{7}+8 n^{6}+\cdots-4733 n^{2}-5130 n+16200\right)}{135 n^{2}(n-1)^{2}(n+1)^{2}(n+2)^{2}} \\
r_{2}(n) & :=\frac{(n-2)(n-3)(n+4)(n+3)\left(7 n^{12}+42 n^{11}+\ldots\right)}{145800 n^{4}(n-1)^{4}(n+1)^{4}(n+2)^{4}}
\end{aligned}
$$

## Further Reading

- Survey article: Creative telescoping for holonomic functions. DOI: 10.1007/978-3-7091-1616-6_7, arXiv:1307.4554.
- PhD thesis: Advanced applications of the holonomic systems approach (RISC, Johannes Kepler University, Linz, Austria).
- Software package: HolonomicFunctions (user's guide). https://risc.jku.at/sw/holonomicfunctions/
- Electromagnetic waves application: Method, device and computer program product for determining an electromagnetic near field of a field excitation source for an electrical system (with J. Schöberl and P. Paule), Patents EP2378444 and US8868382.
- Combinatorial determinants: Binomial determinants for tiling problems yield to the holonomic ansatz (with H . Du, T. Thanatipanonda, E. Wong), DOI: 10.1016/j.ejc.2021.103437.
- 20V determinants: Determinant evaluations inspired by Di Francesco's determinant for twenty-vertex configurations (with C. Krattenthaler and M. Schlosser), arXiv:2401.08481.

